

# ENHANCEMENT IN LEACH PROTOCOL FOR WIRELESS SENSOR NETWORKS

<sup>1</sup>Priyanka Bawa, <sup>2</sup>Bhupendra Singh, <sup>3</sup>Dr.Vinod Shokeen

<sup>1</sup>M.Tech scholar, Department of Electronics and Communication ,Amity University,Noida,India

<sup>2</sup>Assistant Professor-I, Department of Electronics and Communication ,Amity University,Noida,India

<sup>3</sup>Assistant Professor-III, Department of Electronics and Communication ,Amity University,Noida,India

**Abstract**— Wireless Sensor Networks are used for gathering high precision data where the setting up of wired infrastructure is not possible i.e. too difficult or too costly. Wireless Sensor Network consists of large number of sensor nodes with limited communication, computational, sensing capability & irreplaceable power resources. Sensor nodes have constraints like limited power resources and bandwidth.. The main aim is to evenly distribute the energy load among the sensor nodes in order to overcome the problem of overly utilized sensor nodes that will run out of energy as compared to other sensor nodes. In this paper I have proposed an energy-optimised protocol for wireless sensor networks. The proposed protocol is based on clustering of sensor nodes. There are two phases of the protocol i.e. setup phase and steady phase. In this work I have used MatLab 7.9 version for the simulation work.

**Index Terms**—Wireless sensor networks, base station, Clustering, Leach protocol.

## I. INTRODUCTION

With the advancement in the field of highly integrated digital electronics technology and wireless communication, a new class of distributed system called Wireless Sensor Network (WSN) has come into existence. Wireless Sensor Networks are used for gathering high precision data where setting up of wired infrastructure is not possible or too difficult or too costly. WSN consists of large number of sensor nodes with limited communication, computational, sensing capability & irreplaceable power resources. Some applications of WSN include environmental monitoring, military control and communications, forest fire detection, surveillance, traffic monitoring, home applications and control and agricultural applications [1]. A wireless sensor network system usually includes sensor nodes, sink node and management node. A large number of sensor nodes are deployed in the monitored area, constituting a network through the way of self-organization. The data monitored by sensor nodes is transmitted along other nodes one by one, that will reach the sink node after a multi-hop routing and finally reach the management node through the wired and (or) wireless Internet[2]. The energy, the ability of signal process, storage capacity and communication capability of sensor nodes are very limited.

A primary design goal for wireless sensor networks is to use the energy efficiently[3]. Cluster-based routing algorithm has a better energy utilization rate compared with non-cluster routing algorithm[8].

The basic idea of clustering routing[3][7] is to use the information aggregation mechanism in the cluster head to reduce the amount of data transmission, thereby, reduce the energy dissipation in communication and in turn achieve the purpose of saving energy of the sensor nodes. In the clustering routing algorithms for wireless networks, LEACH (low-energy adaptive clustering hierarchy)[4][5] is well-known because it is simple and efficient. LEACH divides the whole network into several clusters, and the run time of network is broken into many rounds. In each round, the nodes in a cluster compete with each other to be cluster head according to a predefined criterion. In LEACH protocol, all the sensor nodes have the same probability to be a cluster head, which makes the nodes in the network consume energy in a relatively balanced way so as to prolong network lifetime.

Some works have been conducted on the performance and enhancement of LEACH protocol[1][6]. This paper proposes an enhanced version of Leach Protocol which aims at prolonging the lifetime of WSN by changing the criterion of selecting cluster heads.

## II. LEACH[8]

LEACH is the first hierarchical protocol in WSN. LEACH is an adaptive clustering routing protocol proposed by Wendi B. Heinzelman. In many later literatures, it has been considered as the benchmark for other protocols. It has some distinctive characteristics like self-reconfiguration, adjustment of communication range according to distance, schedule of data transmission of individual nodes etc. Moreover, unlike most proposed protocols, LEACH has been implemented on actual hardware (MICAz sensor nodes) [7]. It has some assumptions like fixed-base station location, energy constrained homogeneous nodes and predetermined ratio of cluster heads among all nodes. The operation of LEACH is separated into

a series of equal length time spans. In each of these time spans, cluster head selection, cluster formation and scheduling procedures are completed, respectively, at the very beginning. Cluster heads are selected based on a probabilistic value satisfying the condition that those nodes have not played that role previously. Upon receiving broadcasted advertisement messages from a single or multiple cluster heads, a node sends joining declaration to the nearest cluster head. Cluster heads then create a TDMA schedule and notify its member nodes. The following data transmission phase has the larger chunk of each span, which is also divided into a number of equal frames. In each frame, there is a slot for every member node. Member nodes send data to cluster heads at their slot time. The cluster head aggregates the data and send to the base station.

The operation of LEACH is broken up into rounds, where each round begins with a setup phase, when the clusters are organized, followed by a steady-state phase, when data transfers to the base station occur. In order to minimize overhead, the steady-state phase is long compared to the set-up phase. Initially, when clusters are being created, each node decides whether or not to become a cluster-head for the current round. This decision is based on the suggested percentage of cluster heads for the network (determined a priori) and the number of times the node has been a cluster-head so far. This decision is made by the node  $n$  choosing a random number between 0 and 1. If the number is less than a threshold  $T(n)$ , the node becomes a cluster-head for the current round. The threshold is set as:

$$T(n) = \begin{cases} \frac{P}{1-P*(r \bmod \frac{1}{P})} & \text{if } n \in G \\ 0 & \text{otherwise} \end{cases}$$

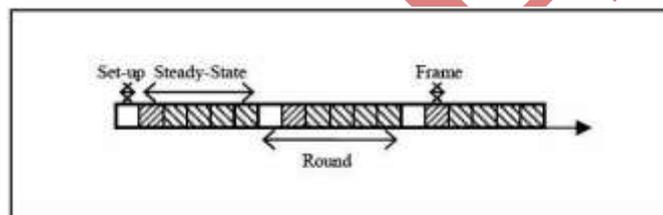


Figure 1:LEACH protocol phases

### III. DESCRIPTION OF THE PROPOSED PROTOCOL

The proposed protocol is based on distributed clustering. In this protocol, all the nodes which are deployed randomly in the required environment to be monitored will be capable of communicating to every other node in the scenario. Every node will advertise its energy levels to all other nodes. The energy level of every node will be checked individually by each node and the nodes which are having energy more than the average energy level of the network are nominated for Cluster Head. The nodes which are at the top ten levels will be selected as Cluster Head initially. These Cluster Heads will advertise their selection to every other node in the network. Those nodes which are not capable of becoming the Cluster Head will wait for the advertisement made by the CH. These nodes will send a join request to the CH which is nearest to them on the basis of received signal strength. The Cluster Head will accept the cluster members and will prepare a TDMA schedule so that each node can communicate to CH without any collision. This TDMA schedule will be sent to the cluster members. Since the nodes will always have data to be send to CH, there are chances that these data may be correlated. The data collected at Cluster Head after the completion of one TDMA Cycle will be fused and aggregated. Now the communication between CH and the Base station will take place with the help of Direct Spread Spectrum so that no collision of data takes place. One by one the CH will be sending the aggregated data to the Base station for monitoring. After one round the process will be repeated again.

### IV. WORK FLOW OF THE PROPOSED PROTOCOL

This section describes diagrammatically the Phases of the proposed protocol. In the Phase 1 i.e. Set Up Phase, every node will advertise its energy level to the other nodes in the network. An average energy level of the network will be calculated in every node. All those nodes which are having energy level greater than or equal to the average energy level will be nominated for Cluster Head. The energy levels will be sorted in descending order by every live node in the network. Since all the nodes are identified by their Identification number, those nodes which are in the top ten level of the energy level will match their identification. After getting selected as cluster head, the CH nodes will advertise their selection and the rest of the non Cluster head nodes will send join request based on the received signal strength of the advertisement of CH nodes to the nearest CH. The Cluster Head will accept their request and will prepare a TDMA slot and the member nodes will send data accordingly. This data will be collected at Cluster Head and it will aggregate that that and will transmit this data to the Base station for monitoring.

### A.SETUP PHASE

The setup phase consists of three parts i.e. Cluster Head Election, Cluster Formation and Schedule Creation. In the Cluster Head Election, nodes having energy level greater than the average energy level of network are nominated and then elected as Cluster Head on the basis of larger energy level. In the Cluster Formation, Non Cluster Head nodes join cluster head based on received signal strength and form the cluster. In the Schedule Creation, the Cluster Head prepares a TDMA schedule for its cluster members for data transmission. A node in the network checks whether it is selected as Cluster Head or not. If the node is not selected as Cluster Head, It will wait for the announcement to be made by the Cluster Head nodes. If the node is selected as Cluster Head, It will advertise its selection to the other nodes in the network so that the Non Cluster Head nodes can join it to become the cluster members. The Cluster Head will wait for the Join Request to be made by the non cluster heads for a time period known a prior. Non cluster head nodes will send Join Request to the Cluster Head which is nearest to them based on the received signal strength of the advertisement made by the Cluster Head nodes. The Cluster Head after getting the responses from the nodes will then prepare a TDMA schedule for its members and will again advertise this schedule only to its cluster members. This schedule will be saved by the cluster members and will send data on the basis of the allotted TDMA slot. After this the steady phase comes into play and after the completion of a round this process will be repeated again. In this way a complete round takes place.

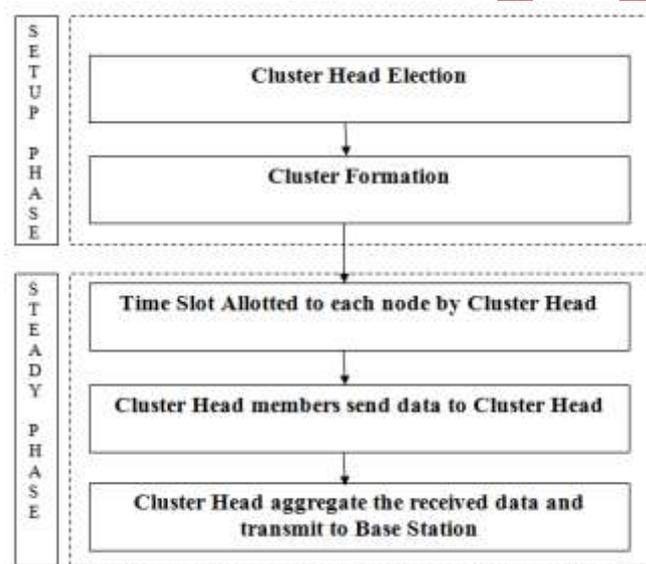


Figure 2: Workflow of the Proposed Protocol

#### 1.Cluster Head Election

In this phase, every node in the network advertises its energy level to the every other node of the network. After this advertisement every node possesses the energy level of each and every node of the network. Now every node will arrange the received energy level of other nodes in the ascending order. The average energy of the network can be calculated as follows:

$$\text{Average Energy} = \frac{\sum_{i=1}^n \text{Energy of each node } i}{\text{Total number of alive nodes}}$$

Those nodes whose energy level is less than the average energy level of the network will be discarded from the Cluster head nomination criteria. Now every node will match its ID with the top ten nodes having the highest energy level. If the ID matches, the node is selected as the Cluster Head. If the ID does not match, the node will wait for the Cluster Head announcements.

#### 2.Cluster Formation

After getting selected as the Cluster Head, now the Cluster Heads will let the other nodes in the network know that they have chosen the role of Cluster Head for the current round. Now the Cluster head will advertise a Join Request

message to all other nodes using a non persistent carrier sense multiple access (CSMA) MAC protocol [9]. This message consists of the Node ID and the Header of this message will be different from normal data messages. Each and every node which is non Cluster Heads will wait for this Join Request Message. The Cluster head will wait for the response for a time period known a prior. Every node will receive this announcement. Now the Non Cluster Head nodes will arrange the received announcements in the decreasing order of the received signal strength. Received Signal Strength is chosen because larger the received signal strength, more the node closer to it unless and until there is an obstacle in between the nodes. The nodes will check which Cluster Head is nearer to it i.e. whose received signal strength is strongest. And will select that Cluster Head. Now the node needs to inform the Cluster Head about its selection. The node will send a Join Response message which will consist of Node ID and a header differentiating it from data messages using a non persistent carrier sense multiple access (CSMA) MAC protocol back to the closest Cluster Head. The Cluster Head will accept the Response and add the node as its cluster member.

### 3. Schedule Creation

Since the Cluster Head also acts as Local data collection center, it will control and coordinate the data transmission between the cluster members. Now all the Cluster Heads will advertise their Cluster member information in between the Cluster Head nodes. The Cluster Head with the largest number of cluster member will set up a time period for the TDMA one cycle. For example, if one of the Cluster Head consists of 20 members and other Cluster Head consist of 10 members then the former Cluster Head will provide one second time allotment for each cluster member and the later Cluster Head will allot two seconds slot for each cluster member. This is done so that the TDMA cycle of each and every Cluster Head should complete at the same time. After the preparation of TDMA schedule by the Cluster Head, the Cluster Head sends this schedule to all its cluster members. Now the cluster members will send their data on the time slot allotted to them. Suppose there are 15 nodes in a cluster. After transmitting the data to the cluster Head, the node will go in sleep mode with its sensor active. Hence the energy of the node can be saved in sleep mode

### B. STEADY PHASE

This phase is the second phase of the proposed protocol. The steady phase consists of Data collection & Aggregation along with Data Transmission

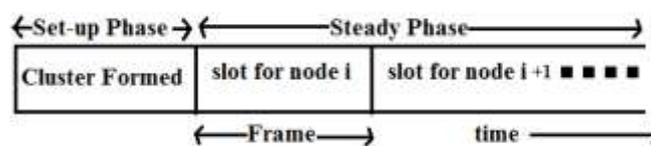


Figure3: steady phase

In Data collection, the cluster head gather data from every cluster member and fuse the data so that no redundancy left. In the data transmission, the Cluster Heads transmit the data one by one to the Base Station for monitoring.

### 1. Data Collection and Aggregation

Data Collection and Aggregation is important part of the Wireless Sensor Network. As the wireless sensor network is bounded by its limited energy supply, therefore emphasis should be made to transmit as less data as possible because transmission consumes most of the power of the wireless sensor network. Aggregation of data is required because if two homogeneous sensors are placed in a same geographical area, then the data sensed by them will be correlated or redundant. Hence Aggregation of data will reduce redundancy thereby decreasing the energy consumption of the network.

The Data collection operation is broken into frames where member nodes send their data to the cluster head at most once per frame during their allotted TDMA slot. The time duration for each slot is fixed where the node sends data to the Cluster Head. The data collected is then fused or aggregated by the Cluster Head so that less number of bits of data is to be sending to the Base Station. The time slot depends on the number of cluster members. It is presumed that the nodes are all time synchronized and the setup phase of the network is initiated at the same time. This is done so that each round of the clusters of the network should get completed at the same time.

### 2. Data Transmission

In this the data collected at Cluster Head is send to the Base Station. To reduce the energy dissipation, each node of the network uses power control mechanism. In this mechanism all nodes in the network can control their transmission and reception power as per the requirement. Suppose the maximum transmission range of a node is 200m, and the distance between the node and its Cluster Head is 35m, now if the node is transmitting the data with its maximum range, then it is dissipating its power uselessly in the environment. So the nodes can adjust the transmission power level as needed i.e. for

above example the node and the Cluster head will be transmitting the data with a maximum transmission power range of 40m so that the energy could be saved.

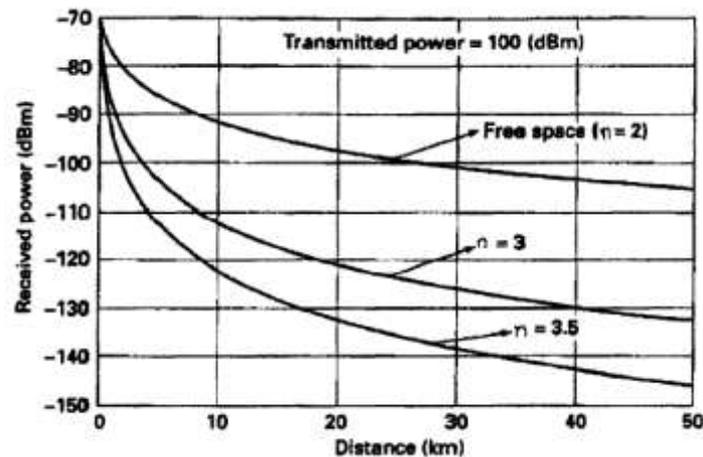


Figure 4:Wireless Communication Radio Model

The data to be sent to Base Station from Cluster Head is done by using fixed spreading code and CSMA. Whenever the Cluster Head has data to be sent to Base Station, it will first sense the channel to ensure that no other Cluster Head is sending the data to the Base Station. If the channel is busy or other Cluster Head are transmitting the data, then the Cluster Head will wait for a time period and then again sense the channel. Now if the channel is free the Cluster Head will send the data using DSSS. DSSS is used for data transmission because the number of cluster heads may vary after some rounds. Other channelization technique like FDMA can be used but it is harder to allot frequency dynamically. The drawback of DSSS is that it requires tight synchronization timing.

### V. ENERGY CONSUMPTION MODEL

Wireless Communication Energy model consists of five major blocks: transmitter’s digital physical layer (PHY) and Medium Access Control (MAC) functions, transmitter’s RF Circuitry including the Phase Locked Loop and Voltage Control Oscillator (VCO), power amplifier, receiver’s RF circuitry and receiver’s digital functions [10].In the proposed protocol, First Order Radio model is assumed [4]. In this model the radio dissipates  $E_{elec} = 50$  nJ/bit to run the transmitter or receiver circuitry and  $E_{amp} = 100$  pJ/bit/m<sup>2</sup> for the transmit amplifier to achieve an acceptable  $E_b/N_o$ . These parameters are slightly better than the current state-of-the-art in radio design

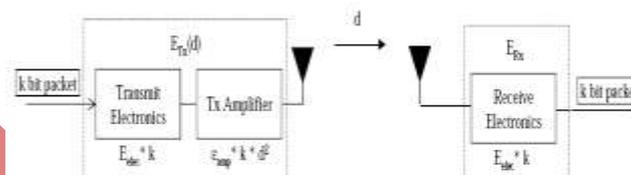


Figure 5:First Order Radio Model[4]

Operation	Energy Dissipated
Transmitter Electronics ( $E_{Tx-elec}$ )	50 nJ/bit
Receiver Electronics ( $E_{Rx-elec}$ ) ( $E_{Tx-elec} = E_{Rx-elec} = E_{elec}$ )	
Transmit Amplifier ( $\epsilon_{amp}$ )	100 pJ/bit/m <sup>2</sup>

Table 1:Radio Model’s Operations energy dissipation[8]

We also assume an r energy loss due to channel transmission. Thus, to transmit a k-bit message a distance d using our radio model, the radio expends:

$$E_{TX}(k,d) = E_{TX-elec}(k) + E_{TX-amp}(k,d)$$

$$E_{TX}(k,d) = E_{elec} * k + \epsilon_{amp} * k * d^2$$

And to receive this message, the radio expends:

$$E_{RX}(k) = E_{RX} - E_{elec}(k)$$

$$E_{RX}(k) = E_{elec} * k$$

For these parameters, receiving a message is also costlier, hence the protocol should try to minimize not only the transmitting distance but also the number of transmit and receive operations. In the proposed protocol, an assumption is made such that the energy dissipated in transmitting a message from a Node A to Node B is equal to the energy dissipated in transmitting a message from Node B to Node A. In the proposed protocol it is also assumed that sensor node has data all the time to transmit to the Cluster Head.

## VI. SIMULATION AND RESULTS

MatLab is used as a simulation platform. With MatLab, it is easier to generate a GUI to show the scenario of the Wireless Sensor Network. There are many functions provided in the MatLab which makes the simulation work easier. MatLab provides quicker configuring the parameters. The different plots can be made in the same simulation with the help of MatLab. Simulation in MatLab is quicker and less complex. There are many protocols which are already implemented in MatLab. Therefore the proposed protocol is implemented in MatLab Version 7.9

### A. Parameters Description

The simulation is carried out in a network grid of 100mX100m. There are two scenarios where the position of the Base station is varied. In one scenario, the location of the Base station is (50, 50) and in another scenario, the Base station is at (175, 50). The numbers of nodes taken are 100. But for the extension of protocol the number of nodes is changed to 200 to find out the proposed protocol behaviour. The percentage of Cluster Head is chosen 10 % of the live nodes in the wireless sensor network. This has been chosen 10% by the experimental analysis as shown in figure 6. The energy levels of each node are also varied in both the scenario. The energy levels are set as 0.25J, 0.50J, and 1J. In some of the experiments the energy level is chosen randomly for every sensor node in the network. It is done so that the proposed protocol should ensure it is more efficient than the LEACH protocol [4]. In some experiments the energy levels are set in multiples of 5 to 20 Joule to know about the lifetime pattern of the network as it will reveal that the network lifetime increases proportionately with the energy of the wireless sensor network. In this model the radio dissipates  $E_{elec} = 50$  nJ/bit to run the transmitter or receiver circuitry and  $E_{amp} = 100$  pJ/bit/m<sup>2</sup> for the transmit amplifier.  $E_{DA} = 5$  nJ/bit/signal is the Energy required for data aggregation.  $\epsilon_{fs}$  is the energy dissipated for free space and  $\epsilon_{mp}$  is the energy dissipated for multi path fading in wireless sensor network. Different topologies are generated randomly for simulation. Each simulation result shows the average of independent experiments.

Parameters	Values
Network Grid	From (0,0) to (100,100)
Number of Nodes	100
Base Station Location	(175,50)m / (50,50)m
Initial Energy	0.25 J/ 0.50J/1 J
$\epsilon_{fs}$	10 pJ/bit/m <sup>4</sup>
$\epsilon_{mp}$	0.0013pJ/bit/m <sup>4</sup>
$E_{elec}$	50 nJ/bit
$E_{DA}$	5 nJ/bit/signal

Table 2: Simulation Parameters for Proposed Protocol

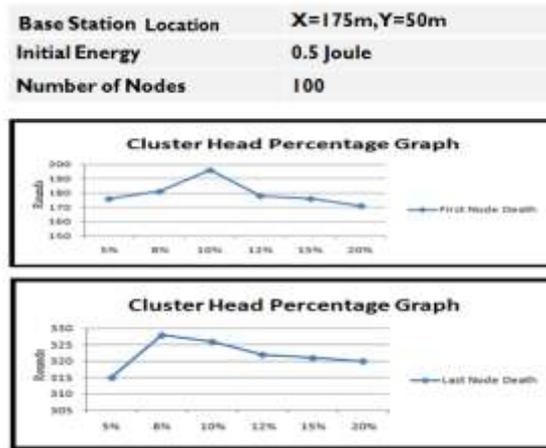


Figure 6: Cluster Head Percentage Selection

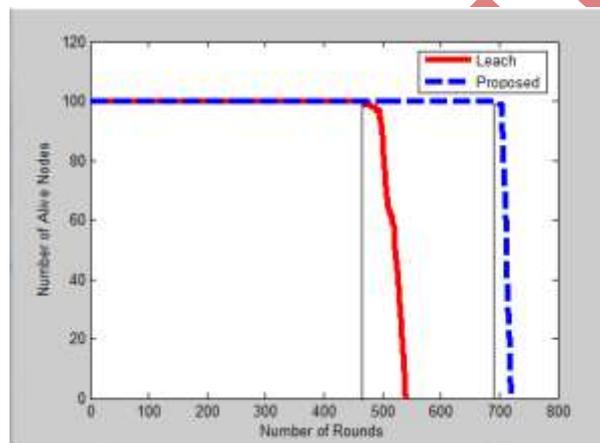


Figure 7: Plot of Alive nodes vs. Round (x=175m, y=50m,  $E_0=0.5J$ )

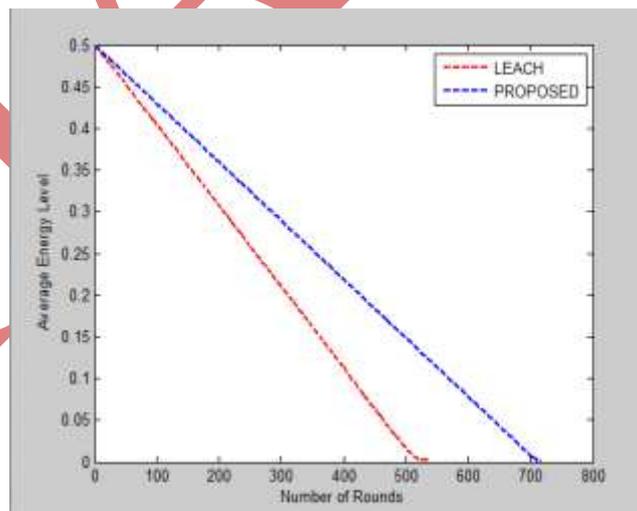


Figure 8: Plot of Average Energy Level of nodes vs. Round (x=175m, y=50m,  $E_0=0.5J$ )

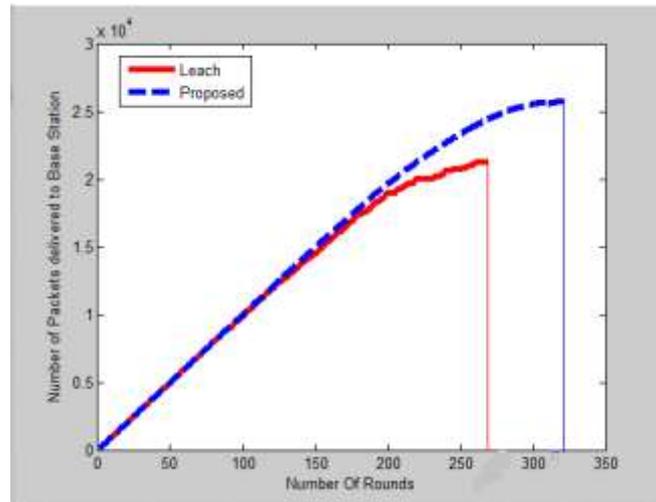


Figure 9: Plot of Packet delivered to BS vs. Round (x=175m, y=50m,  $E_0=0.5J$ )

## VII.CONCLUSION

The proposed protocol selects ten percent of the live node as cluster head for the cluster formation. These cluster head possesses the highest energy level in the wireless sensor network. This is done so because the cluster head utilises more energy as compared to the other nodes in the network as it receives the data packet send by its cluster member, it aggregates and fuses the gathered data and send the data to the Base Station.

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